

igpm

## Topic 6: QR-Iteration with Shift

In this exercise we want to enhance the QR-Iteration from the previous exercise by inclusion of a clever shift-strategy. The *QR-Iteration with Shift* shall then be applied for symmetric and non-symmetric matrices. Since this exercise requires a successful implementation of the QR iteration, it is mandatory to complete the previous exercise.

From now on let  $A \in \mathbb{R}^{n \times n}$  be diagonalisable with eigenvalues  $\lambda_i$ . We assume that all eigenvalues are of distinct magnitude and w.l.o.g. ordered:

$$|\lambda_1| > |\lambda_2| > \ldots > |\lambda_n|$$

The speed of convergence of the QR-iteration depends on the ratios  $|\lambda_{i+1}/\lambda_i|$  of successive eigenvalues. If the ratio is close to 0 then it converges quickly, if it is close to 1 then it slows down considerably.

Similarly to the trick of the shifted inverse iteration, one can here as well introduce a Shift and thereby speed up the convergence — at least for one of the eigenvalues at a time. Let  $\mu \in \mathbb{R}$  be an approximation of the eigenvalue  $\lambda_i$  such that

$$|\mu - \lambda_i| \ll |\mu - \lambda_j| \quad \forall j \neq i.$$

The QR-Iteration applied to the matrix  $A - \mu I$  will then quickly converge for the first eigenvector corresponding to the eigenvalue  $\lambda_i$ :

## Algorithm 1 QR-Iteration with Shift

1:  $A_1 := A, Q_1 := I$ 2: for i = 1, 2, ... do 3: Determine  $\mu$ 4: Compute a QR-decomposition  $(A_i - \mu I) =: \hat{Q}_{i+1}R_{i+1}$ 5:  $A_{i+1} := R_{i+1}\hat{Q}_{i+1} + \mu I$ 6:  $Q_{i+1} := Q_i\hat{Q}_{i+1}$ 7: end for

Algorithm 1 has the following properties:

- $A_i$  converges to an upper triangular matrix, but in general  $A_i \neq Q_i^T A Q_i =: \hat{A}_i$ .
- The diagonal entries of  $\hat{A}_i$  are approximations of the eigenvalues of A.
- The shift parameter  $\mu$  can be chosen differently in each step (without increasing the computational complexity)

In the following we assume that the matrix A is already given in upper Hessenberg form. The shift parameter  $\mu$  in Algorithm 1 shall now be chosen in such a way that it approximates the same eigenvalue in every step (but we do not know the eigenvalues, how should we do this ?). After a

few steps the matrix  $M := \hat{A}_i$  should be of the form

for some  $k \in \{1, \ldots, n\}$ . If the entry  $m_{k+1,k}$  is close enough to 0, then the problem can be split into two subproblems for the matrices  $M_{11} \in \mathbb{R}^{k \times k}$  and  $M_{22} \in \mathbb{R}^{(n-k) \times (n-k)}$ .

 $M_{11}$  and  $M_{22}$  are both of upper Hessenberg form. The algorithm can thus be applied recursively for  $M_{11}$  and  $M_{22}$ . From the solutions of the subproblems we obtain the solution for M. We apply Algorithm 1 until the subproblems are of size  $1 \times 1$  and thus trivial. An important criterion is to decide whether or not entry  $m_{k+1,k}$  in (1) is small enough. In the following the problem is subdivided into subproblems if for given  $\varepsilon > 0$ 

$$|m_{k+1,k}| < \varepsilon. \tag{2}$$

Copy all files from ../../vorlagen/EWP/ex6new into your directory EWP

cp -r ../../vorlagen/EWP/ex6new .

Complete in the file qr\_shift.cc the function

which is supposed to perform the QR-iteration with shift. Exploit the recursive structure in your program. The input matrix A shall not be changed by your procedure. The result  $A_i$  shall be stored in  $A_i$  and the orthogonal factor  $Q_i$  in Q. The parameter eps controls the splitting criterion (2).

Compile your program by make.

**Test:** Start your programm  $qr_shift$  and test it for the matrices A und B. Analyse the output (on the screen).